

**NAME**

`poll`, `ppoll` – wait for some event on a file descriptor

**SYNOPSIS**

```
#include <poll.h>

int poll(struct pollfd *fds, nfd_t nfd, int timeout);

#define _GNU_SOURCE /* See feature_test_macros(7) */
#include <signal.h>
#include <poll.h>

int ppoll(struct pollfd *fds, nfd_t nfd,
          const struct timespec *tmo_p, const sigset_t *sigmask);
```

**DESCRIPTION**

`poll()` performs a similar task to `select(2)`: it waits for one of a set of file descriptors to become ready to perform I/O.

The set of file descriptors to be monitored is specified in the `fds` argument, which is an array of structures of the following form:

```
struct pollfd {
    int    fd;          /* file descriptor */
    short  events;     /* requested events */
    short  revents;    /* returned events */
};
```

The caller should specify the number of items in the `fds` array in `nfd`.

The field `fd` contains a file descriptor for an open file. If this field is negative, then the corresponding `events` field is ignored and the `revents` field returns zero. (This provides an easy way of ignoring a file descriptor for a single `poll()` call: simply negate the `fd` field. Note, however, that this technique can't be used to ignore file descriptor 0.)

The field `events` is an input parameter, a bit mask specifying the events the application is interested in for the file descriptor `fd`. This field may be specified as zero, in which case the only events that can be returned in `revents` are **POLLHUP**, **POLLERR**, and **POLLNVAL** (see below).

The field `revents` is an output parameter, filled by the kernel with the events that actually occurred. The bits returned in `revents` can include any of those specified in `events`, or one of the values **POLLERR**, **POLLHUP**, or **POLLNVAL**. (These three bits are meaningless in the `events` field, and will be set in the `revents` field whenever the corresponding condition is true.)

If none of the events requested (and no error) has occurred for any of the file descriptors, then `poll()` blocks until one of the events occurs.

The `timeout` argument specifies the number of milliseconds that `poll()` should block waiting for a file descriptor to become ready. The call will block until either:

- \* a file descriptor becomes ready;
- \* the call is interrupted by a signal handler; or
- \* the timeout expires.

Note that the `timeout` interval will be rounded up to the system clock granularity, and kernel scheduling delays mean that the blocking interval may overrun by a small amount. Specifying a negative value in `timeout` means an infinite timeout. Specifying a `timeout` of zero causes `poll()` to return immediately, even if no file descriptors are ready.

The bits that may be set/returned in `events` and `revents` are defined in `<poll.h>`:

**POLLIN**

There is data to read.

**POLLPRI**

There is some exceptional condition on the file descriptor. Possibilities include:

- \* There is out-of-band data on a TCP socket (see **tcp(7)**).
- \* A pseudoterminal master in packet mode has seen a state change on the slave (see **ioctl\_tty(2)**).
- \* A *cgroup.events* file has been modified (see **cgroups(7)**).

**POLLOUT**

Writing is now possible, though a write larger than the available space in a socket or pipe will still block (unless **O\_NONBLOCK** is set).

**POLLRDHUP** (since Linux 2.6.17)

Stream socket peer closed connection, or shut down writing half of connection. The **\_GNU\_SOURCE** feature test macro must be defined (before including *any* header files) in order to obtain this definition.

**POLLERR**

Error condition (only returned in *revents*; ignored in *events*). This bit is also set for a file descriptor referring to the write end of a pipe when the read end has been closed.

**POLLHUP**

Hang up (only returned in *revents*; ignored in *events*). Note that when reading from a channel such as a pipe or a stream socket, this event merely indicates that the peer closed its end of the channel. Subsequent reads from the channel will return 0 (end of file) only after all outstanding data in the channel has been consumed.

**POLLNVAL**

Invalid request: *fd* not open (only returned in *revents*; ignored in *events*).

When compiling with **\_XOPEN\_SOURCE** defined, one also has the following, which convey no further information beyond the bits listed above:

**POLLRDNORM**

Equivalent to **POLLIN**.

**POLLRDBAND**

Priority band data can be read (generally unused on Linux).

**POLLWRNORM**

Equivalent to **POLLOUT**.

**POLLWRBAND**

Priority data may be written.

Linux also knows about, but does not use **POLLMSG**.

**ppoll()**

The relationship between **poll()** and **ppoll()** is analogous to the relationship between **select(2)** and **pselect(2)**: like **pselect(2)**, **ppoll()** allows an application to safely wait until either a file descriptor becomes ready or until a signal is caught.

Other than the difference in the precision of the *timeout* argument, the following **ppoll()** call:

```
ready = ppoll(&fds, nfd, tmo_p, &sigmask);
```

is nearly equivalent to *atomically* executing the following calls:

```
sigset_t origmask;
int timeout;
```

```

timeout = (tmo_p == NULL) ? -1 :
          (tmo_p->tv_sec * 1000 + tmo_p->tv_nsec / 1000000);
pthread_sigmask(SIG_SETMASK, &sigmask, &origmask);
ready = poll(&fds, nfds, timeout);
pthread_sigmask(SIG_SETMASK, &origmask, NULL);

```

The above code segment is described as *nearly* equivalent because whereas a negative *timeout* value for **poll()** is interpreted as an infinite timeout, a negative value expressed in *\*tmo\_p* results in an error from **ppoll()**.

See the description of **pselect(2)** for an explanation of why **ppoll()** is necessary.

If the *sigmask* argument is specified as NULL, then no signal mask manipulation is performed (and thus **ppoll()** differs from **poll()** only in the precision of the *timeout* argument).

The *tmo\_p* argument specifies an upper limit on the amount of time that **ppoll()** will block. This argument is a pointer to a structure of the following form:

```

struct timespec {
    long    tv_sec;           /* seconds */
    long    tv_nsec;        /* nanoseconds */
};

```

If *tmo\_p* is specified as NULL, then **ppoll()** can block indefinitely.

## RETURN VALUE

On success, a positive number is returned; this is the number of structures which have nonzero *revents* fields (in other words, those descriptors with events or errors reported). A value of 0 indicates that the call timed out and no file descriptors were ready. On error, -1 is returned, and *errno* is set appropriately.

## ERRORS

### EFAULT

The array given as argument was not contained in the calling program's address space.

### EINTR

A signal occurred before any requested event; see **signal(7)**.

### EINVAL

The *nfds* value exceeds the **RLIMIT\_NOFILE** value.

### EINVAL

(**ppoll()**) The timeout value expressed in *\*ip* is invalid (negative).

### ENOMEM

There was no space to allocate file descriptor tables.

## VERSIONS

The **poll()** system call was introduced in Linux 2.1.23. On older kernels that lack this system call, the glibc (and the old Linux libc) **poll()** wrapper function provides emulation using **select(2)**.

The **ppoll()** system call was added to Linux in kernel 2.6.16. The **ppoll()** library call was added in glibc 2.4.

## CONFORMING TO

**poll()** conforms to POSIX.1-2001 and POSIX.1-2008. **ppoll()** is Linux-specific.

## NOTES

The operation of **poll()** and **ppoll()** is not affected by the **O\_NONBLOCK** flag.

On some other UNIX systems, **poll()** can fail with the error **EAGAIN** if the system fails to allocate kernel-internal resources, rather than **ENOMEM** as Linux does. POSIX permits this behavior. Portable programs may wish to check for **EAGAIN** and loop, just as with **EINTR**.

Some implementations define the nonstandard constant **INFTIM** with the value -1 for use as a *timeout* for **poll()**. This constant is not provided in glibc.

For a discussion of what may happen if a file descriptor being monitored by **poll()** is closed in another thread, see **select(2)**.

### C library/kernel differences

The Linux **ppoll()** system call modifies its *tmo\_p* argument. However, the glibc wrapper function hides this behavior by using a local variable for the timeout argument that is passed to the system call. Thus, the glibc **ppoll()** function does not modify its *tmo\_p* argument.

The raw **ppoll()** system call has a fifth argument, *size\_t sigsetsize*, which specifies the size in bytes of the *sigmask* argument. The glibc **ppoll()** wrapper function specifies this argument as a fixed value (equal to *sizeof(kernel\_sigset\_t)*). See **sigprocmask(2)** for a discussion on the differences between the kernel and the libc notion of the sigset.

### BUGS

See the discussion of spurious readiness notifications under the BUGS section of **select(2)**.

### SEE ALSO

**restart\_syscall(2)**, **select(2)**, **select\_tut(2)**, **epoll(7)**, **time(7)**

### COLOPHON

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