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### ***Rocky Enterprise Linux 9.2 Manual Pages on command 'semtimeop.2'***

#### ***\$ man semtimeop.2***

SEMOP(2)                   Linux Programmer's Manual                   SEMOP(2)

#### NAME

semop, semtimeop - System V semaphore operations

#### SYNOPSIS

```
#include <sys/types.h>
#include <sys/ipc.h>
#include <sys/sem.h>

int semop(int semid, struct sembuf *sops, size_t nsops);

int semtimeop(int semid, struct sembuf *sops, size_t nsops,
              const struct timespec *timeout);
```

Feature Test Macro Requirements for glibc (see feature\_test\_macros(7)):

```
semtimeop(): _GNU_SOURCE
```

#### DESCRIPTION

Each semaphore in a System V semaphore set has the following associated values:

```
unsigned short semval; /* semaphore value */
unsigned short semzcnt; /* # waiting for zero */
unsigned short semncnt; /* # waiting for increase */
```

```
pid_t      sempid; /* PID of process that last
```

semop() performs operations on selected semaphores in the set indicated by semid. Each of the nsops elements in the array pointed to by sops is a structure that specifies an operation to be performed on a single semaphore. The elements of this structure are of type struct sembuf, containing the following members:

```
unsigned short sem_num; /* semaphore number */  
short         sem_op; /* semaphore operation */  
short         sem_flg; /* operation flags */
```

Flags recognized in sem\_flg are IPC\_NOWAIT and SEM\_UNDO. If an operation specifies SEM\_UNDO, it will be automatically undone when the process terminates.

The set of operations contained in sops is performed in array order, and atomically, that is, the operations are performed either as a complete unit, or not at all. The behavior of the system call if not all operations can be performed immediately depends on the presence of the IPC\_NOWAIT flag in the individual sem\_flg fields, as noted below.

Each operation is performed on the sem\_num-th semaphore of the semaphore set, where the first semaphore of the set is numbered 0. There are three types of operation, distinguished by the value of sem\_op.

If sem\_op is a positive integer, the operation adds this value to the semaphore value (semval). Furthermore, if SEM\_UNDO is specified for this operation, the system subtracts the value sem\_op from the semaphore adjustment (semadj) value for this semaphore. This operation can always proceed; it never forces a thread to wait. The calling process must have alter permission on the semaphore set.

If sem\_op is zero, the process must have read permission on the semaphore set. This is a "wait-for-zero" operation: if semval is zero, the operation can immediately proceed. Otherwise, if IPC\_NOWAIT is specified in sem\_flg, semop() fails with errno set to EAGAIN (and none of the operations in sops is performed). Otherwise, semzcnt (the count of threads waiting until this semaphore's value becomes zero) is incremented by one and the thread sleeps until one of the following occurs:

? semval becomes 0, at which time the value of semzcnt is decremented.

? The semaphore set is removed: semop() fails, with errno set to EIDRM.

? The calling thread catches a signal: the value of semzcnt is decremented and semop() fails, with errno set to EINTR.

If sem\_op is less than zero, the process must have alter permission on the semaphore set. If semval is greater than or equal to the absolute value of sem\_op, the operation can proceed immediately: the absolute value of sem\_op is subtracted from semval, and, if SEM\_UNDO is specified for this operation, the system adds the absolute value of sem\_op to the semaphore adjustment (semadj) value for this semaphore. If the absolute value of sem\_op is greater than semval, and IPC\_NOWAIT is specified in sem\_flg, semop() fails, with errno set to EAGAIN (and none of the operations in sops is performed). Otherwise, semncnt (the counter of threads waiting for this semaphore's value to increase) is incremented by one and the thread sleeps until one of the following occurs:

? semval becomes greater than or equal to the absolute value of sem\_op: the operation now proceeds, as described above.

? The semaphore set is removed from the system: semop() fails, with errno set to EIDRM.

? The calling thread catches a signal: the value of semncnt is decremented and semop() fails, with errno set to EINTR.

On successful completion, the sempid value for each semaphore specified in the array pointed to by sops is set to the caller's process ID. In addition, the sem\_otime is set to the current time.

#### semtimedop()

semtimedop() behaves identically to semop() except that in those cases where the calling thread would sleep, the duration of that sleep is limited by the amount of elapsed time specified by the timespec structure whose address is passed in the timeout argument. (This sleep interval will be rounded up to the system clock granularity, and kernel scheduling delays mean that the interval may overrun by a small amount.) If the specified time limit has been reached, semtimedop()

fails with `errno` set to `EAGAIN` (and none of the operations in `sops` is performed). If the `timeout` argument is `NULL`, then `semtimedop()` behaves exactly like `semop()`.

Note that if `semtimedop()` is interrupted by a signal, causing the call to fail with the error `EINTR`, the contents of `timeout` are left unchanged.

## RETURN VALUE

If successful, `semop()` and `semtimedop()` return 0; otherwise they return -1 with `errno` indicating the error.

## ERRORS

On failure, `errno` is set to one of the following:

**E2BIG** The argument `nsops` is greater than `SEMOPM`, the maximum number of operations allowed per system call.

**EACCES** The calling process does not have the permissions required to perform the specified semaphore operations, and does not have the `CAP_IPC_OWNER` capability in the user namespace that governs its IPC namespace.

**EAGAIN** An operation could not proceed immediately and either `IPC_NOWAIT` was specified in `sem_flg` or the time limit specified in `timeout` expired.

**EFAULT** An address specified in either the `sops` or the `timeout` argument isn't accessible.

**EFBIG** For some operation the value of `sem_num` is less than 0 or greater than or equal to the number of semaphores in the set.

**EIDRM** The semaphore set was removed.

**EINTR** While blocked in this system call, the thread caught a signal; see `signal(7)`.

**EINVAL** The semaphore set doesn't exist, or `semid` is less than zero, or `nsops` has a nonpositive value.

**ENOMEM** The `sem_flg` of some operation specified `SEM_UNDO` and the system does not have enough memory to allocate the undo structure.

**ERANGE** For some operation `sem_op+semval` is greater than `SEMVMX`, the implementation dependent maximum value for `semval`.

## VERSIONS

semtimeop() first appeared in Linux 2.5.52, and was subsequently backported into kernel 2.4.22. Glibc support for semtimeop() first appeared in version 2.3.3.

## CONFORMING TO

POSIX.1-2001, POSIX.1-2008, SVr4.

## NOTES

The inclusion of `<sys/types.h>` and `<sys/ipc.h>` isn't required on Linux or by any version of POSIX. However, some old implementations required the inclusion of these header files, and the SVID also documented their inclusion. Applications intended to be portable to such old systems may need to include these header files.

The `sem_undo` structures of a process aren't inherited by the child produced by `fork(2)`, but they are inherited across an `execve(2)` system call.

`semop()` is never automatically restarted after being interrupted by a signal handler, regardless of the setting of the `SA_RESTART` flag when establishing a signal handler.

A semaphore adjustment (`semadj`) value is a per-process, per-semaphore integer that is the negated sum of all operations performed on a semaphore specifying the `SEM_UNDO` flag. Each process has a list of `semadj` values—one value for each semaphore on which it has operated using `SEM_UNDO`. When a process terminates, each of its per-semaphore `semadj` values is added to the corresponding semaphore, thus undoing the effect of that process's operations on the semaphore (but see `BUGS` below).

When a semaphore's value is directly set using the `SETVAL` or `SETALL` request to `semctl(2)`, the corresponding `semadj` values in all processes are cleared. The `clone(2)` `CLONE_SYSVSEM` flag allows more than one process to share a `semadj` list; see `clone(2)` for details.

The `semval`, `sempid`, `semzcnt`, and `semnct` values for a semaphore can all be retrieved using appropriate `semctl(2)` calls.

## Semaphore limits

The following limits on semaphore set resources affect the `semop()`

call:

SEMOPM Maximum number of operations allowed for one semop() call. Be?

fore Linux 3.19, the default value for this limit was 32. Since Linux 3.19, the default value is 500. On Linux, this limit can be read and modified via the third field of /proc/sys/kernel/sem. Note: this limit should not be raised above 1000, because of the risk of that semop() fails due to kernel memory fragmentation when allocating memory to copy the sops array.

SEMVMX Maximum allowable value for semval: implementation dependent (32767).

The implementation has no intrinsic limits for the adjustment on exit maximum value (SEMAEM), the system wide maximum number of undo structures (SEMMNU) and the per-process maximum number of undo entries system parameters.

## BUGS

When a process terminates, its set of associated semadj structures is used to undo the effect of all of the semaphore operations it performed with the SEM\_UNDO flag. This raises a difficulty: if one (or more) of these semaphore adjustments would result in an attempt to decrease a semaphore's value below zero, what should an implementation do? One possible approach would be to block until all the semaphore adjustments could be performed. This is however undesirable since it could force process termination to block for arbitrarily long periods. Another possibility is that such semaphore adjustments could be ignored altogether (somewhat analogously to failing when IPC\_NOWAIT is specified for a semaphore operation). Linux adopts a third approach: decreasing the semaphore value as far as possible (i.e., to zero) and allowing process termination to proceed immediately.

In kernels 2.6.x, x <= 10, there is a bug that in some circumstances prevents a thread that is waiting for a semaphore value to become zero from being woken up when the value does actually become zero. This bug is fixed in kernel 2.6.11.

## EXAMPLES

The following code segment uses `semop()` to atomically wait for the value of semaphore 0 to become zero, and then increment the semaphore value by one.

```
struct sembuf sops[2];
int semid;
/* Code to set semid omitted */
sops[0].sem_num = 0;    /* Operate on semaphore 0 */
sops[0].sem_op = 0;    /* Wait for value to equal 0 */
sops[0].sem_flg = 0;
sops[1].sem_num = 0;    /* Operate on semaphore 0 */
sops[1].sem_op = 1;    /* Increment value by one */
sops[1].sem_flg = 0;
if (semop(semid, sops, 2) == -1) {
    perror("semop");
    exit(EXIT_FAILURE);
}
```

A further example of the use of `semop()` can be found in `shmop(2)`.

#### SEE ALSO

`clone(2)`, `semctl(2)`, `semget(2)`, `sigaction(2)`, `capabilities(7)`,  
`sem_overview(7)`, `sysvipc(7)`, `time(7)`

#### COLOPHON

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