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Rocky Enterprise Linux 9.2 Manual Pages on command 'getauxval.3'

# \$ man getauxval.3

GETAUXVAL(3)

Linux Programmer's Manual

GETAUXVAL(3)

# NAME

getauxval - retrieve a value from the auxiliary vector

# SYNOPSIS

#include <sys/auxv.h>

unsigned long getauxval(unsigned long type);

# DESCRIPTION

The getauxval() function retrieves values from the auxiliary vector, a

mechanism that the kernel's ELF binary loader uses to pass certain in?

formation to user space when a program is executed.

Each entry in the auxiliary vector consists of a pair of values: a type

that identifies what this entry represents, and a value for that type.

Given the argument type, getauxval() returns the corresponding value.

The value returned for each type is given in the following list. Not

all type values are present on all architectures.

# AT\_BASE

The base address of the program interpreter (usually, the dy? namic linker).

### AT\_BASE\_PLATFORM

A pointer to a string (PowerPC and MIPS only). On PowerPC, this

identifies the real platform; may differ from AT\_PLATFORM. On

MIPS, this identifies the ISA level (since Linux 5.7).

### AT\_CLKTCK

The frequency with which times(2) counts. This value can also

be obtained via sysconf(\_SC\_CLK\_TCK).

## AT\_DCACHEBSIZE

The data cache block size.

## AT\_EGID

The effective group ID of the thread.

## AT\_ENTRY

The entry address of the executable.

## AT\_EUID

The effective user ID of the thread.

## AT\_EXECFD

File descriptor of program.

### AT\_EXECFN

A pointer to a string containing the pathname used to execute

the program.

# AT\_FLAGS

Flags (unused).

# AT\_FPUCW

Used FPU control word (SuperH architecture only). This gives

some information about the FPU initialization performed by the

kernel.

# AT\_GID The real group ID of the thread.

### AT\_HWCAP

An architecture and ABI dependent bit-mask whose settings indi? cate detailed processor capabilities. The contents of the bit mask are hardware dependent (for example, see the kernel source file arch/x86/include/asm/cpufeature.h for details relating to the Intel x86 architecture; the value returned is the first 32-bit word of the array described there). A human-readable version of the same information is available via /proc/cpuinfo.

## AT\_HWCAP2 (since glibc 2.18)

Further machine-dependent hints about processor capabilities.

AT\_ICACHEBSIZE

The instruction cache block size.

## AT\_L1D\_CACHEGEOMETRY

Geometry of the L1 data cache, encoded with the cache line size

in bytes in the bottom 16 bits and the cache associativity in

the next 16 bits. The associativity is such that if N is the

16-bit value, the cache is N-way set associative.

## AT\_L1D\_CACHESIZE

The L1 data cache size.

## AT\_L1I\_CACHEGEOMETRY

Geometry of the L1 instruction cache, encoded as for

AT\_L1D\_CACHEGEOMETRY.

### AT\_L1I\_CACHESIZE

The L1 instruction cache size.

### AT\_L2\_CACHEGEOMETRY

Geometry of the L2 cache, encoded as for AT\_L1D\_CACHEGEOMETRY.

### AT\_L2\_CACHESIZE

The L2 cache size.

### AT\_L3\_CACHEGEOMETRY

Geometry of the L3 cache, encoded as for AT\_L1D\_CACHEGEOMETRY.

### AT\_L3\_CACHESIZE

The L3 cache size.

### AT\_PAGESZ

The system page size (the same value returned by sysconf(\_SC\_PA?

GESIZE)).

### AT\_PHDR

The address of the program headers of the executable.

### AT\_PHENT

The size of program header entry.

#### AT\_PHNUM

The number of program headers.

### AT\_PLATFORM

A pointer to a string that identifies the hardware platform that the program is running on. The dynamic linker uses this in the interpretation of rpath values.

## AT\_RANDOM

The address of sixteen bytes containing a random value.

### AT\_SECURE

Has a nonzero value if this executable should be treated se? curely. Most commonly, a nonzero value indicates that the process is executing a set-user-ID or set-group-ID binary (so that its real and effective UIDs or GIDs differ from one an? other), or that it gained capabilities by executing a binary file that has capabilities (see capabilities(7)). Alterna? tively, a nonzero value may be triggered by a Linux Security Module. When this value is nonzero, the dynamic linker disables the use of certain environment variables (see Id-linux.so(8)) and glibc changes other aspects of its behavior. (See also se? cure\_getenv(3).)

### AT\_SYSINFO

The entry point to the system call function in the vDSO. Not present/needed on all architectures (e.g., absent on x86-64).

### AT\_SYSINFO\_EHDR

The address of a page containing the virtual Dynamic Shared Ob? ject (vDSO) that the kernel creates in order to provide fast im?

plementations of certain system calls.

### AT\_UCACHEBSIZE

The unified cache block size.

AT\_UID The real user ID of the thread.

#### **RETURN VALUE**

On success, getauxval() returns the value corresponding to type. If

type is not found, 0 is returned.

## ERRORS

ENOENT (since glibc 2.19)

No entry corresponding to type could be found in the auxiliary vector.

#### VERSIONS

The getauxval() function was added to glibc in version 2.16.

#### ATTRIBUTES

For an explanation of the terms used in this section, see at?

tributes(7).

?Interface ? Attribute ? Value ?

?getauxval() ? Thread safety ? MT-Safe ?

#### CONFORMING TO

This function is a nonstandard glibc extension.

#### NOTES

The primary consumer of the information in the auxiliary vector is the dynamic linker, Id-linux.so(8). The auxiliary vector is a convenient and efficient shortcut that allows the kernel to communicate a certain set of standard information that the dynamic linker usually or always needs. In some cases, the same information could be obtained by system calls, but using the auxiliary vector is cheaper. The auxiliary vector resides just above the argument list and environ?

ment in the process address space. The auxiliary vector supplied to a program can be viewed by setting the LD\_SHOW\_AUXV environment variable when running a program:

\$ LD\_SHOW\_AUXV=1 sleep 1

The auxiliary vector of any process can (subject to file permissions) be obtained via /proc/[pid]/auxv; see proc(5) for more information.

#### BUGS

Before the addition of the ENOENT error in glibc 2.19, there was no way to unambiguously distinguish the case where type could not be found from the case where the value corresponding to type was zero.

# SEE ALSO

secure\_getenv(3), vdso(7), ld-linux.so(8)

# COLOPHON

This page is part of release 5.10 of the Linux man-pages project. A description of the project, information about reporting bugs, and the latest version of this page, can be found at https://www.kernel.org/doc/man-pages/.

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