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## **Red Hat Enterprise Linux Release 9.2 Manual Pages on 'pthread\_atfork.3' command**

**\$ man pthread\_atfork.3**

PTHREAD\_ATFORK(3)      Linux Programmer's Manual      PTHREAD\_ATFORK(3)

### NAME

pthread\_atfork - register fork handlers

### SYNOPSIS

```
#include <pthread.h>

int pthread_atfork(void (*prepare)(void), void (*parent)(void),
                  void (*child)(void));
```

Link with -pthread.

### DESCRIPTION

The pthread\_atfork() function registers fork handlers that are to be executed when fork(2) is called by this thread. The handlers are executed in the context of the thread that calls fork(2).

Three kinds of handler can be registered:

- \* prepare specifies a handler that is executed before fork(2) processing starts.
- \* parent specifies a handler that is executed in the parent process after fork(2) processing completes.
- \* child specifies a handler that is executed in the child process after fork(2) processing completes.

Any of the three arguments may be NULL if no handler is needed in the corresponding phase of fork(2) processing.

### RETURN VALUE

On success, pthread\_atfork() returns zero. On error, it returns an error

ror number. `pthread_atfork()` may be called multiple times by a thread, to register multiple handlers for each phase. The handlers for each phase are called in a specified order: the prepare handlers are called in reverse order of registration; the parent and child handlers are called in the order of registration.

## ERRORS

`ENOMEM` Could not allocate memory to record the form handler entry.

## CONFORMING TO

POSIX.1-2001, POSIX.1-2008.

## NOTES

When `fork(2)` is called in a multithreaded process, only the calling thread is duplicated in the child process. The original intention of `pthread_atfork()` was to allow the calling thread to be returned to a consistent state. For example, at the time of the call to `fork(2)`, other threads may have locked mutexes that are visible in the user-space memory duplicated in the child. Such mutexes would never be unlocked, since the threads that placed the locks are not duplicated in the child. The intent of `pthread_atfork()` was to provide a mechanism whereby the application (or a library) could ensure that mutexes and other process and thread state would be restored to a consistent state. In practice, this task is generally too difficult to be practicable.

After a `fork(2)` in a multithreaded process returns in the child, the child should call only async-signal-safe functions (see `signal-safety(7)`) until such time as it calls `execve(2)` to execute a new program.

POSIX.1 specifies that `pthread_atfork()` shall not fail with the error `EINTR`.

## SEE ALSO

`fork(2)`, `atexit(3)`, `threads(7)`

## COLOPHON

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