



Red Hat Enterprise Linux Release 9.2 Manual Pages on 'fts_open.3' command

\$ man fts_open.3

FTS(3) Linux Programmer's Manual FTS(3)

NAME

fts, fts_open, fts_read, fts_children, fts_set, fts_close - traverse a file hierarchy

SYNOPSIS

```
#include <sys/types.h>
#include <sys/stat.h>
#include <fts.h>

FTS *fts_open(char * const *path_argv, int options,
              int (*compar)(const FTSENT **, const FTSENT **));

FTSENT *fts_read(FTS *ftsp);

FTSENT *fts_children(FTS *ftsp, int instr);

int fts_set(FTS *ftsp, FTSENT *f, int instr);

int fts_close(FTS *ftsp);
```

DESCRIPTION

The fts functions are provided for traversing file hierarchies. A simple overview is that the fts_open() function returns a "handle" (of type FTS *) that refers to a file hierarchy "stream". This handle is then supplied to the other fts functions. The function fts_read() returns a pointer to a structure describing one of the files in the file hierarchy. The function fts_children() returns a pointer to a linked list of structures, each of which describes one of the files contained in a directory in the hierarchy.

In general, directories are visited two distinguishable times; in preorder (before any of their descendants are visited) and in postorder (after all of their descendants have been visited). Files are visited once. It is possible to walk the hierarchy "logically" (visiting the files that symbolic links point to) or physically (visiting the symbolic links themselves), order the walk of the hierarchy or prune and/or revisit portions of the hierarchy.

Two structures (and associated types) are defined in the include file `<fts.h>`. The first type is `FTS`, the structure that represents the file hierarchy itself. The second type is `FTSENT`, the structure that represents a file in the file hierarchy. Normally, an `FTSENT` structure is returned for every file in the file hierarchy. In this manual page, "file" and "FTSENT structure" are generally interchangeable.

The `FTSENT` structure contains fields describing a file. The structure contains at least the following fields (there are additional fields that should be considered private to the implementation):

```
typedef struct _ftsent {
    unsigned short fts_info; /* flags for FTSENT structure */
    char *fts_accpath; /* access path */
    char *fts_path; /* root path */
    short fts_pathlen; /* strlen(fts_path) +
                       strlen(fts_name) */
    char *fts_name; /* filename */
    short fts_namelen; /* strlen(fts_name) */
    short fts_level; /* depth (-1 to N) */
    int fts_errno; /* file errno */
    long fts_number; /* local numeric value */
    void *fts_pointer; /* local address value */
    struct _ftsent *fts_parent; /* parent directory */
    struct _ftsent *fts_link; /* next file structure */
    struct _ftsent *fts_cycle; /* cycle structure */
    struct stat *fts_statp; /* stat(2) information */
} FTSENT;
```

These fields are defined as follows:

fts_info

One of the following values describing the returned FTSENT structure and the file it represents. With the exception of directories without errors (FTS_D), all of these entries are terminal, that is, they will not be revisited, nor will any of their descendants be visited.

FTS_D A directory being visited in preorder.

FTS_DC A directory that causes a cycle in the tree. (The fts_cycle field of the FTSENT structure will be filled in as well.)

FTS_DEFAULT

Any FTSENT structure that represents a file type not explicitly described by one of the other fts_info values.

FTS_DNR

A directory which cannot be read. This is an error return, and the fts_errno field will be set to indicate what caused the error.

FTS_DOT

A file named "." or ".." which was not specified as a filename to fts_open() (see FTS_SEEDOT).

FTS_DP A directory being visited in postorder. The contents of the FTSENT structure will be unchanged from when it was returned in preorder, that is, with the fts_info field set to FTS_D.

FTS_ERR

This is an error return, and the fts_errno field will be set to indicate what caused the error.

FTS_F A regular file.

FTS_NS A file for which no stat(2) information was available.

The contents of the fts_statp field are undefined. This is an error return, and the fts_errno field will be set to indicate what caused the error.

FTS_NSOK

A file for which no stat(2) information was requested.

The contents of the fts_statp field are undefined.

FTS_SL A symbolic link.

FTS_SLNONE

A symbolic link with a nonexistent target. The contents of the fts_statp field reference the file characteristic information for the symbolic link itself.

fts_accpath

A path for accessing the file from the current directory.

fts_path

The path for the file relative to the root of the traversal.

This path contains the path specified to fts_open() as a prefix.

fts_pathlen

The sum of the lengths of the strings referenced by fts_path and fts_name.

fts_name

The name of the file.

fts_namelen

The length of the string referenced by fts_name.

fts_level

The depth of the traversal, numbered from -1 to N, where this file was found. The FTSENT structure representing the parent of the starting point (or root) of the traversal is numbered -1, and the FTSENT structure for the root itself is numbered 0.

fts_errno

If fts_children() or fts_read() returns an FTSENT structure whose fts_info field is set to FTS_DNR, FTS_ERR, or FTS_NS, the fts_errno field contains the error number (i.e., the errno value) specifying the cause of the error. Otherwise, the contents of the fts_errno field are undefined.

fts_number

This field is provided for the use of the application program

and is not modified by the fts functions. It is initialized to 0.

fts_pointer

This field is provided for the use of the application program and is not modified by the fts functions. It is initialized to NULL.

fts_parent

A pointer to the FTSENT structure referencing the file in the hierarchy immediately above the current file, that is, the directory of which this file is a member. A parent structure for the initial entry point is provided as well, however, only the fts_level, fts_number, and fts_pointer fields are guaranteed to be initialized.

fts_link

Upon return from the fts_children() function, the fts_link field points to the next structure in the NULL-terminated linked list of directory members. Otherwise, the contents of the fts_link field are undefined.

fts_cycle

If a directory causes a cycle in the hierarchy (see FTS_DC), either because of a hard link between two directories, or a symbolic link pointing to a directory, the fts_cycle field of the structure will point to the FTSENT structure in the hierarchy that references the same file as the current FTSENT structure. Otherwise, the contents of the fts_cycle field are undefined.

fts_statp

A pointer to stat(2) information for the file.

A single buffer is used for all of the paths of all of the files in the file hierarchy. Therefore, the fts_path and fts_accpath fields are guaranteed to be null-terminated only for the file most recently returned by fts_read(). To use these fields to reference any files represented by other FTSENT structures will require that the path buffer be modified using the information contained in that FTSENT structure's

fts_pathlen field. Any such modifications should be undone before further calls to fts_read() are attempted. The fts_name field is always null-terminated.

fts_open()

The fts_open() function takes a pointer to an array of character pointers naming one or more paths which make up a logical file hierarchy to be traversed. The array must be terminated by a null pointer.

There are a number of options, at least one of which (either FTS_LOGICAL or FTS_PHYSICAL) must be specified. The options are selected by ORing the following values:

FTS_COMFOLLOW

This option causes any symbolic link specified as a root path to be followed immediately whether or not FTS_LOGICAL is also specified.

FTS_LOGICAL

This option causes the fts routines to return FTSENT structures for the targets of symbolic links instead of the symbolic links themselves. If this option is set, the only symbolic links for which FTSENT structures are returned to the application are those referencing nonexistent files. Either FTS_LOGICAL or FTS_PHYSICAL must be provided to the fts_open() function.

FTS_NOCHDIR

As a performance optimization, the fts functions change directories as they walk the file hierarchy. This has the side-effect that an application cannot rely on being in any particular directory during the traversal. The FTS_NOCHDIR option turns off this optimization, and the fts functions will not change the current directory. Note that applications should not themselves change their current directory and try to access files unless FTS_NOCHDIR is specified and absolute pathnames were provided as arguments to fts_open().

FTS_NOSTAT

By default, returned FTSENT structures reference file character?

istic information (the statp field) for each file visited. This option relaxes that requirement as a performance optimization, allowing the fts functions to set the fts_info field to FTS_NSOK and leave the contents of the statp field undefined.

FTS_PHYSICAL

This option causes the fts routines to return FTSENT structures for symbolic links themselves instead of the target files they point to. If this option is set, FTSENT structures for all symbolic links in the hierarchy are returned to the application. Either FTS_LOGICAL or FTS_PHYSICAL must be provided to the fts_open() function.

FTS_SEEDOT

By default, unless they are specified as path arguments to fts_open(), any files named "." or ".." encountered in the file hierarchy are ignored. This option causes the fts routines to return FTSENT structures for them.

FTS_XDEV

This option prevents fts from descending into directories that have a different device number than the file from which the descent began.

The argument compar() specifies a user-defined function which may be used to order the traversal of the hierarchy. It takes two pointers to pointers to FTSENT structures as arguments and should return a negative value, zero, or a positive value to indicate if the file referenced by its first argument comes before, in any order with respect to, or after, the file referenced by its second argument. The fts_accpath, fts_path, and fts_pathlen fields of the FTSENT structures may never be used in this comparison. If the fts_info field is set to FTS_NS or FTS_NSOK, the fts_statp field may not either. If the compar() argument is NULL, the directory traversal order is in the order listed in path_argv for the root paths, and in the order listed in the directory for everything else.

The `fts_read()` function returns a pointer to an `FTSENT` structure describing a file in the hierarchy. Directories (that are readable and do not cause cycles) are visited at least twice, once in preorder and once in postorder. All other files are visited at least once. (Hard links between directories that do not cause cycles or symbolic links to symbolic links may cause files to be visited more than once, or directories more than twice.)

If all the members of the hierarchy have been returned, `fts_read()` returns `NULL` and sets the external variable `errno` to 0. If an error unrelated to a file in the hierarchy occurs, `fts_read()` returns `NULL` and sets `errno` appropriately. If an error related to a returned file occurs, a pointer to an `FTSENT` structure is returned, and `errno` may or may not have been set (see `fts_info`).

The `FTSENT` structures returned by `fts_read()` may be overwritten after a call to `fts_close()` on the same file hierarchy stream, or, after a call to `fts_read()` on the same file hierarchy stream unless they represent a file of type directory, in which case they will not be overwritten until after a call to `fts_read()` after the `FTSENT` structure has been returned by the function `fts_read()` in postorder.

`fts_children()`

The `fts_children()` function returns a pointer to an `FTSENT` structure describing the first entry in a `NULL`-terminated linked list of the files in the directory represented by the `FTSENT` structure most recently returned by `fts_read()`. The list is linked through the `fts_link` field of the `FTSENT` structure, and is ordered by the user-specified comparison function, if any. Repeated calls to `fts_children()` will recreate this linked list.

As a special case, if `fts_read()` has not yet been called for a hierarchy, `fts_children()` will return a pointer to the files in the logical directory specified to `fts_open()`, that is, the arguments specified to `fts_open()`. Otherwise, if the `FTSENT` structure most recently returned by `fts_read()` is not a directory being visited in preorder, or the directory does not contain any files, `fts_children()` returns `NULL` and

sets `errno` to zero. If an error occurs, `fts_children()` returns `NULL` and sets `errno` appropriately.

The `FTSENT` structures returned by `fts_children()` may be overwritten after a call to `fts_children()`, `fts_close()`, or `fts_read()` on the same file hierarchy stream.

The `instr` argument is either zero or the following value:

`FTS_NAMEONLY`

Only the names of the files are needed. The contents of all the fields in the returned linked list of structures are undefined with the exception of the `fts_name` and `fts_namelen` fields.

`fts_set()`

The function `fts_set()` allows the user application to determine further processing for the file `f` of the stream `ftsp`. The `fts_set()` function returns 0 on success, and -1 if an error occurs.

The `instr` argument is either 0 (meaning "do nothing") or one of the following values:

`FTS_AGAIN`

Revisit the file; any file type may be revisited. The next call to `fts_read()` will return the referenced file. The `fts_stat` and `fts_info` fields of the structure will be reinitialized at that time, but no other fields will have been changed. This option is meaningful only for the most recently returned file from `fts_read()`. Normal use is for postorder directory visits, where it causes the directory to be revisited (in both preorder and postorder) as well as all of its descendants.

`FTS_FOLLOW`

The referenced file must be a symbolic link. If the referenced file is the one most recently returned by `fts_read()`, the next call to `fts_read()` returns the file with the `fts_info` and `fts_statp` fields reinitialized to reflect the target of the symbolic link instead of the symbolic link itself. If the file is one of those most recently returned by `fts_children()`, the `fts_info` and `fts_statp` fields of the structure, when returned by

`fts_read()`, will reflect the target of the symbolic link instead of the symbolic link itself. In either case, if the target of the symbolic link does not exist, the fields of the returned structure will be unchanged and the `fts_info` field will be set to `FTS_SLNONE`.

If the target of the link is a directory, the preorder return, followed by the return of all of its descendants, followed by a postorder return, is done.

FTS_SKIP

No descendants of this file are visited. The file may be one of those most recently returned by either `fts_children()` or `fts_read()`.

`fts_close()`

The `fts_close()` function closes the file hierarchy stream referred to by `ftsp` and restores the current directory to the directory from which `fts_open()` was called to open `ftsp`. The `fts_close()` function returns 0 on success, and -1 if an error occurs.

ERRORS

The function `fts_open()` may fail and set `errno` for any of the errors specified for `open(2)` and `malloc(3)`.

The function `fts_close()` may fail and set `errno` for any of the errors specified for `chdir(2)` and `close(2)`.

The functions `fts_read()` and `fts_children()` may fail and set `errno` for any of the errors specified for `chdir(2)`, `malloc(3)`, `opendir(3)`, `read(2)`, `readdir(3)`, and `stat(2)`.

In addition, `fts_children()`, `fts_open()`, and `fts_set()` may fail and set `errno` as follows:

`EINVAL` options or `instr` was invalid.

VERSIONS

These functions are available in Linux since `glibc2`.

ATTRIBUTES

For an explanation of the terms used in this section, see `attributes(7)`.

