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Red Hat Enterprise Linux Release 9.2 Manual Pages on 'calloc.3' command

\$ man calloc.3

MALLOC(3) Linux Programmer's Manual MALLOC(3)

NAME

malloc, free, calloc, realloc, reallocarray - allocate and free dynamic memory

SYNOPSIS

```
#include <stdlib.h>

void *malloc(size_t size);

void free(void *ptr);

void *calloc(size_t nmemb, size_t size);

void *realloc(void *ptr, size_t size);

void *reallocarray(void *ptr, size_t nmemb, size_t size);
```

Feature Test Macro Requirements for glibc (see feature_test_macros(7)):

reallocarray():

Since glibc 2.29:

 _DEFAULT_SOURCE

Glibc 2.28 and earlier:

 _GNU_SOURCE

DESCRIPTION

The malloc() function allocates size bytes and returns a pointer to the allocated memory. The memory is not initialized. If size is 0, then malloc() returns either NULL, or a unique pointer value that can later be successfully passed to free().

The free() function frees the memory space pointed to by ptr, which

must have been returned by a previous call to `malloc()`, `calloc()`, or `realloc()`. Otherwise, or if `free(ptr)` has already been called before, undefined behavior occurs. If `ptr` is `NULL`, no operation is performed. The `calloc()` function allocates memory for an array of `nmemb` elements of `size` bytes each and returns a pointer to the allocated memory. The memory is set to zero. If `nmemb` or `size` is 0, then `calloc()` returns either `NULL`, or a unique pointer value that can later be successfully passed to `free()`. If the multiplication of `nmemb` and `size` would result in integer overflow, then `calloc()` returns an error. By contrast, an integer overflow would not be detected in the following call to `malloc()`, with the result that an incorrectly sized block of memory would be allocated:

```
malloc(nmemb * size);
```

The `realloc()` function changes the size of the memory block pointed to by `ptr` to `size` bytes. The contents will be unchanged in the range from the start of the region up to the minimum of the old and new sizes. If the new size is larger than the old size, the added memory will not be initialized. If `ptr` is `NULL`, then the call is equivalent to `malloc(size)`, for all values of `size`; if `size` is equal to zero, and `ptr` is not `NULL`, then the call is equivalent to `free(ptr)`. Unless `ptr` is `NULL`, it must have been returned by an earlier call to `malloc()`, `calloc()`, or `realloc()`. If the area pointed to was moved, a `free(ptr)` is done.

The `reallocarray()` function changes the size of the memory block pointed to by `ptr` to be large enough for an array of `nmemb` elements, each of which is `size` bytes. It is equivalent to the call

```
realloc(ptr, nmemb * size);
```

However, unlike that `realloc()` call, `reallocarray()` fails safely in the case where the multiplication would overflow. If such an overflow occurs, `reallocarray()` returns `NULL`, sets `errno` to `ENOMEM`, and leaves the original block of memory unchanged.

RETURN VALUE

The `malloc()` and `calloc()` functions return a pointer to the allocated

memory, which is suitably aligned for any built-in type. On error, these functions return NULL. NULL may also be returned by a successful call to malloc() with a size of zero, or by a successful call to calloc() with nmemb or size equal to zero.

The free() function returns no value.

The realloc() function returns a pointer to the newly allocated memory, which is suitably aligned for any built-in type, or NULL if the request failed. The returned pointer may be the same as ptr if the allocation was not moved (e.g., there was room to expand the allocation in-place), or different from ptr if the allocation was moved to a new address. If size was equal to 0, either NULL or a pointer suitable to be passed to free() is returned. If realloc() fails, the original block is left untouched; it is not freed or moved.

On success, the reallocarray() function returns a pointer to the newly allocated memory. On failure, it returns NULL and the original block of memory is left untouched.

ERRORS

calloc(), malloc(), realloc(), and reallocarray() can fail with the following error:

ENOMEM Out of memory. Possibly, the application hit the RLIMIT_AS or RLIMIT_DATA limit described in getrlimit(2).

VERSIONS

reallocarray() first appeared in glibc in version 2.26.

ATTRIBUTES

For an explanation of the terms used in this section, see attributes(7).

??

?Interface ? Attribute ? Value ?

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?malloc(), free(), ? Thread safety ? MT-Safe ?

?calloc(), realloc() ? ? ?

??

CONFORMING TO

malloc(), free(), calloc(), realloc(): POSIX.1-2001, POSIX.1-2008, C89, C99.

reallocarray() is a nonstandard extension that first appeared in OpenBSD 5.6 and FreeBSD 11.0.

NOTES

By default, Linux follows an optimistic memory allocation strategy.

This means that when malloc() returns non-NULL there is no guarantee that the memory really is available. In case it turns out that the system is out of memory, one or more processes will be killed by the OOM killer. For more information, see the description of /proc/sys/vm/overcommit_memory and /proc/sys/vm/oom_adj in proc(5), and the Linux kernel source file Documentation/vm/overcommit-accounting.rst.

Normally, malloc() allocates memory from the heap, and adjusts the size of the heap as required, using sbrk(2). When allocating blocks of memory larger than MMAP_THRESHOLD bytes, the glibc malloc() implementation allocates the memory as a private anonymous mapping using mmap(2). MMAP_THRESHOLD is 128 kB by default, but is adjustable using malloc(3). Prior to Linux 4.7 allocations performed using mmap(2) were unaffected by the RLIMIT_DATA resource limit; since Linux 4.7, this limit is also enforced for allocations performed using mmap(2).

To avoid corruption in multithreaded applications, mutexes are used internally to protect the memory-management data structures employed by these functions. In a multithreaded application in which threads simultaneously allocate and free memory, there could be contention for these mutexes. To scalably handle memory allocation in multithreaded applications, glibc creates additional memory allocation arenas if mutex contention is detected. Each arena is a large region of memory that is internally allocated by the system (using brk(2) or mmap(2)), and managed with its own mutexes.

SUSv2 requires malloc(), calloc(), and realloc() to set errno to ENOMEM upon failure. Glibc assumes that this is done (and the glibc versions of these routines do this); if you use a private malloc implementation

that does not set `errno`, then certain library routines may fail without having a reason in `errno`.

Crashes in `malloc()`, `calloc()`, `realloc()`, or `free()` are almost always related to heap corruption, such as overflowing an allocated chunk or freeing the same pointer twice.

The `malloc()` implementation is tunable via environment variables; see `mallopt(3)` for details.

SEE ALSO

`valgrind(1)`, `brk(2)`, `mmap(2)`, `alloca(3)`, `malloc_get_state(3)`,
`malloc_info(3)`, `malloc_trim(3)`, `malloc_usable_size(3)`, `mallopt(3)`,
`mcheck(3)`, `mtrace(3)`, `posix_memalign(3)`

For details of the GNU C library implementation, see

[?https://sourceware.org/glibc/wiki/MallocInternals?](https://sourceware.org/glibc/wiki/MallocInternals?).

COLOPHON

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